

Thread Tailor

Dynamically Weaving Threads Together for
Efficient, Adaptive Parallel Applications

Janghaeng Lee, Haicheng Wu,
Madhumitha Ravichandran, Nathan Clark

**Georgia
Tech**



College of
Computing

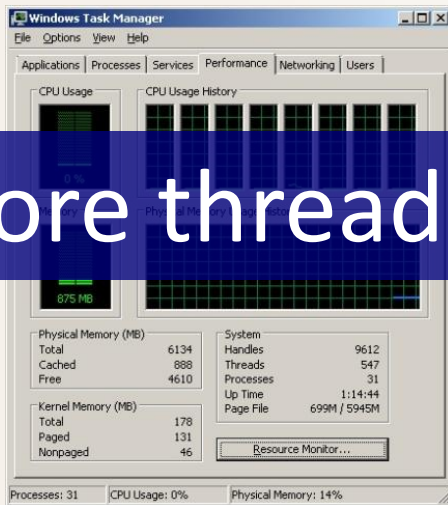


Motivation

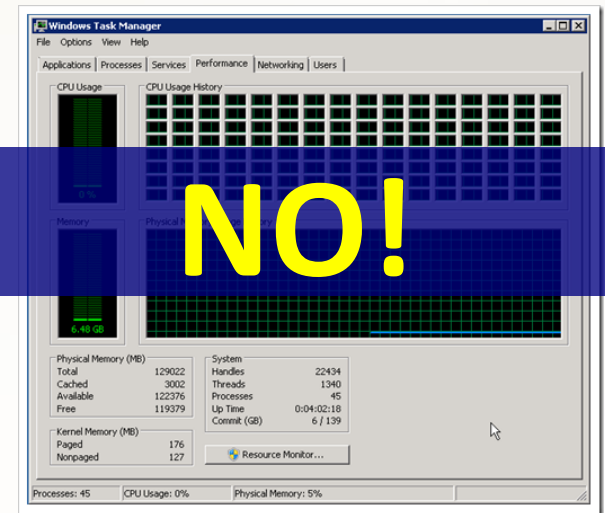
- Hardware Trends
 - Put more cores in a single chip

More threads always win? →

NO!



2009



201X

- CPU intensive programs
 - Exploits Thread Level Parallelism



Optimal Number of Threads

- Too many threads
 - More synchronization
 - More contention for system resources
- Too few threads
 - Resource underutilization
- Who can decide the number?
 - **Not a programmer**



Why NOT?

- Input changes
 - Various working-set size

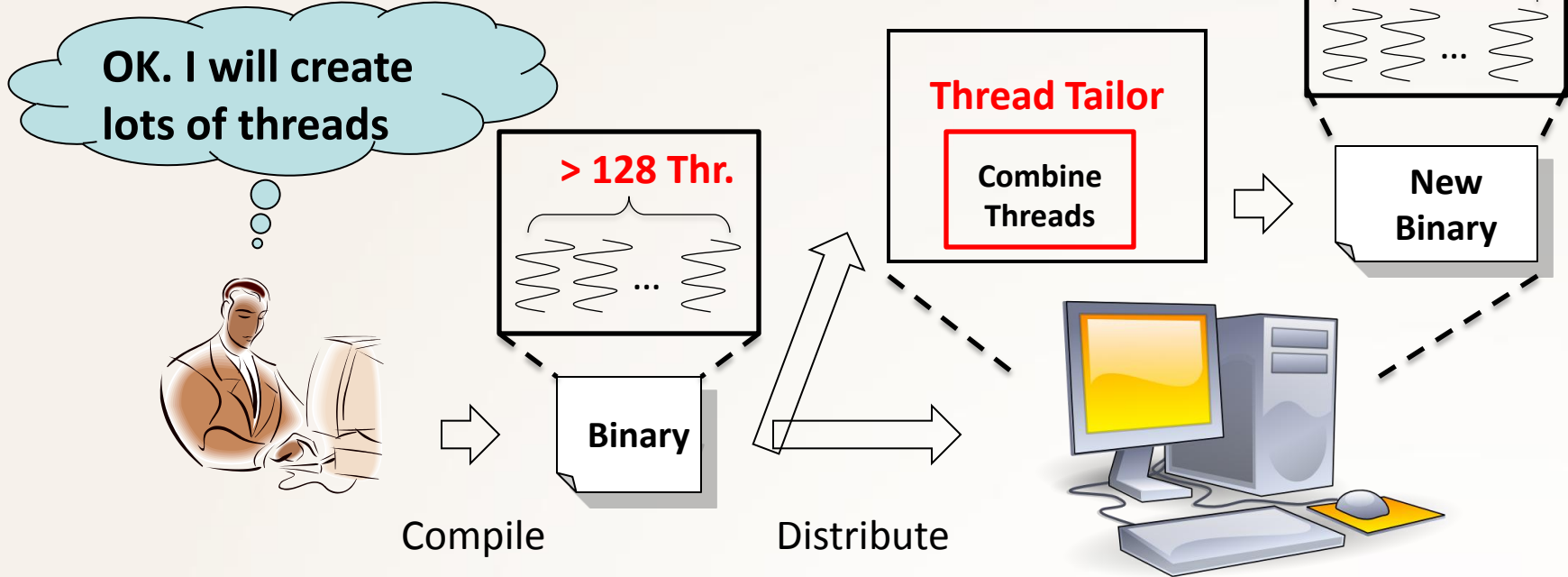
- The system changes

Decision must be made at runtime

- Various available resources

- Hardware changes
 - Various L2/L3 cache structure / size, etc.

Proposal

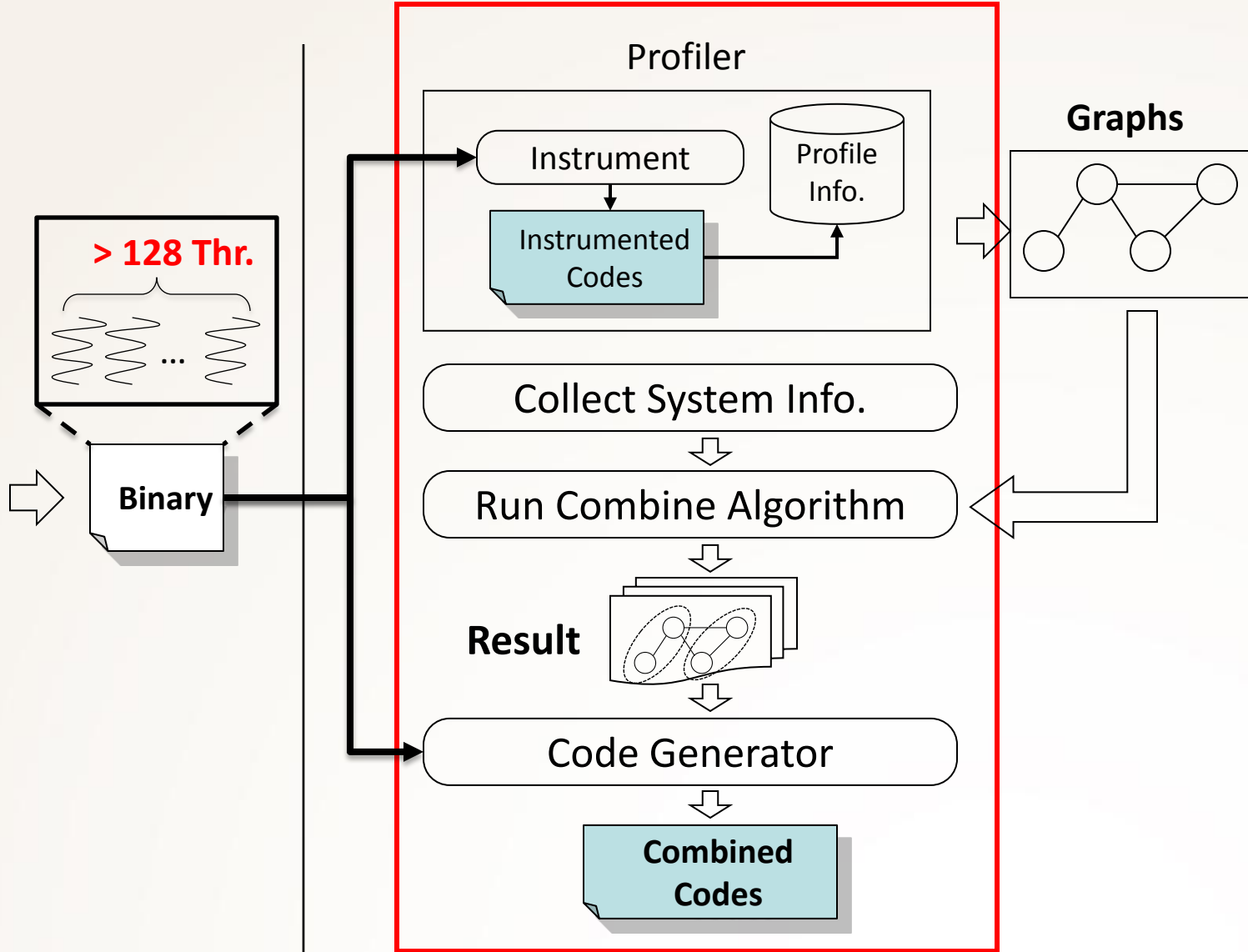


- Combining Threads

- Group Several Threads into a Single Thread

- Threads in the same group are executed in serial
- Executed on the SAME core

Details

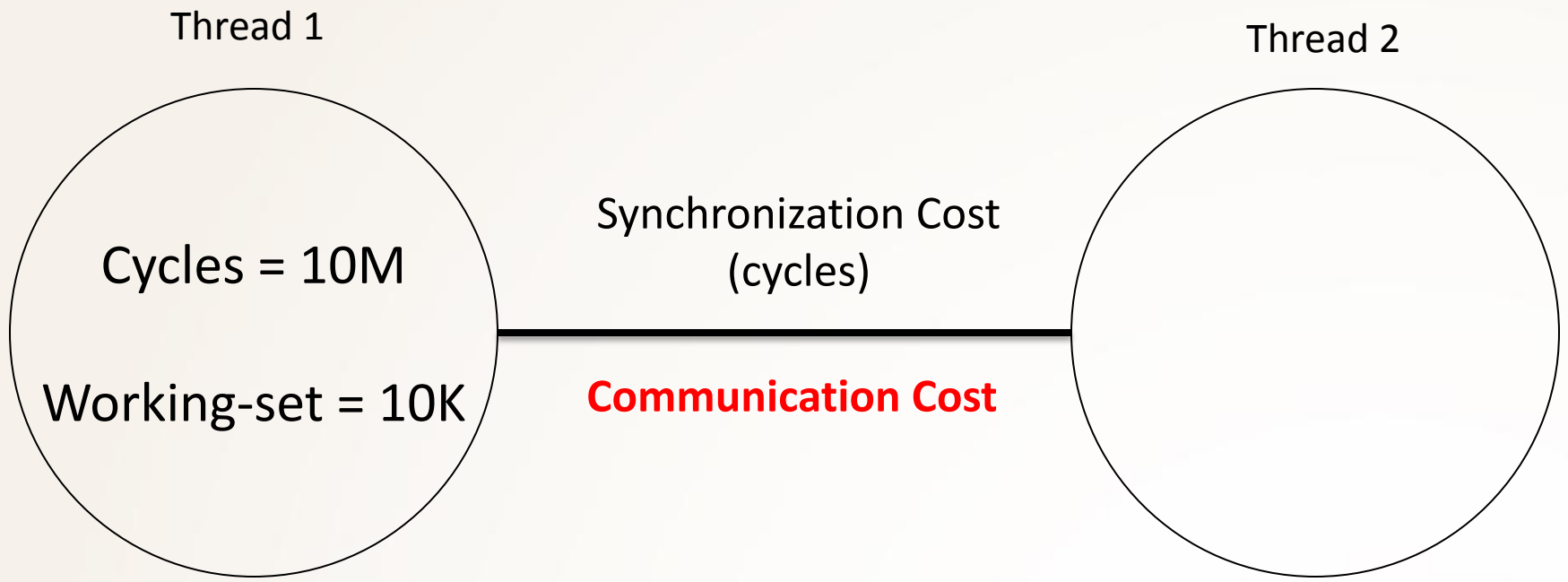


Development | Distribution

Thread Tailor



Graph Construction





Communication Cost

- Intuition : STORE Instruction causes coherence miss in cache
- Log Memory Access per Thread

Thread 1		
Address	LD Count	ST Count
...
0x00001234	5	10
0x00001338	4	9
...
0x00004000	7	7
...

LD LD

ST ST

Graph

① — 29 — ②

Thread 2		
Address	LD Count	ST Count
...
0x00001234	0	7
0x00002000	4	4
...
0x00004000	3	8
...

$$0x00001234: \text{MIN}(5, 7) + \text{MIN}(10, 0) + \text{MIN}(10, 7) = 12$$

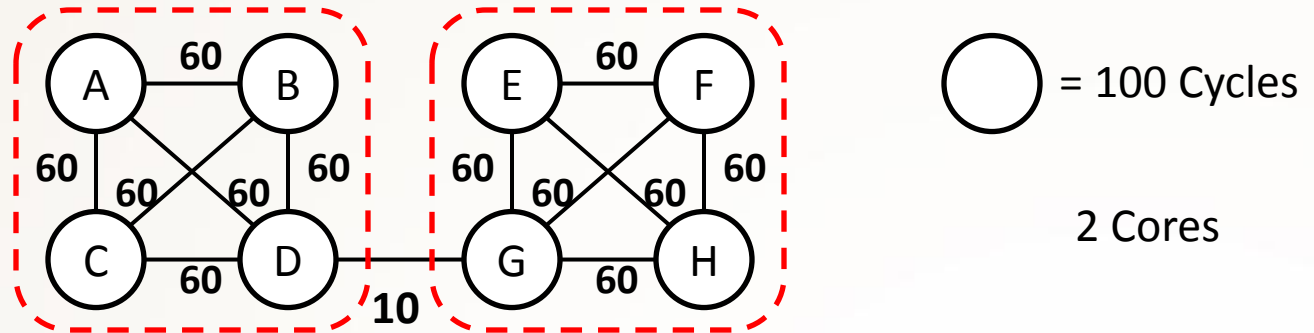
$$0x00004000: \text{MIN}(7, 8) + \text{MIN}(7, 3) + \text{MIN}(7, 8) = 17$$

$$\text{Total Communication Cost: } 12 + 17 = \mathbf{29}$$



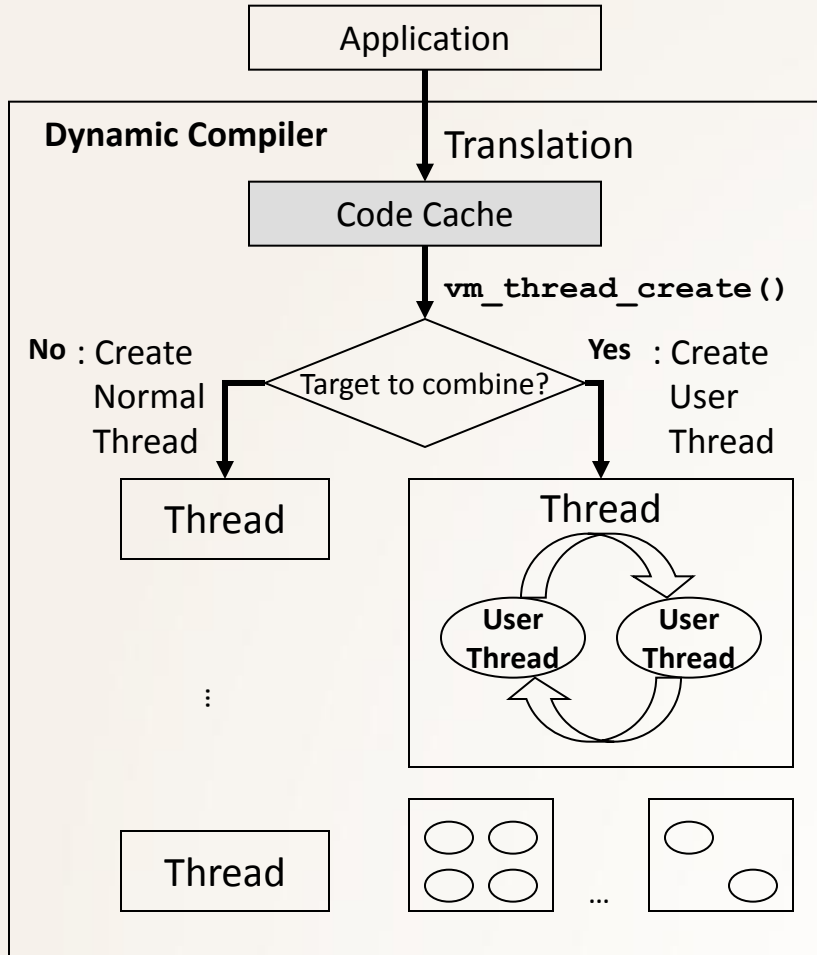
Combining Algorithm

- Kernighan-Lin(KL) Graph Partitioning Heuristic
 - Goal : Minimize Execution Cycles
 - Precondition : Combined Threads \leq Cores



Partition 1	Partition 2	Partition 1 Cycle Estimation	Partition2 Cycle Estimation	Move From	Move Node
B C D G	A E F H	210	220	2	A
A B C D G	E F H	130	120	1	G
A B C D	E F G H	40	40	1	D

Thread Combining



Replace Thread APIs with Wrapper Functions

Wrapper Function for Thread Creation

Context Switched by Dynamic Compiler

Serially Execute User Threads in Real Thread



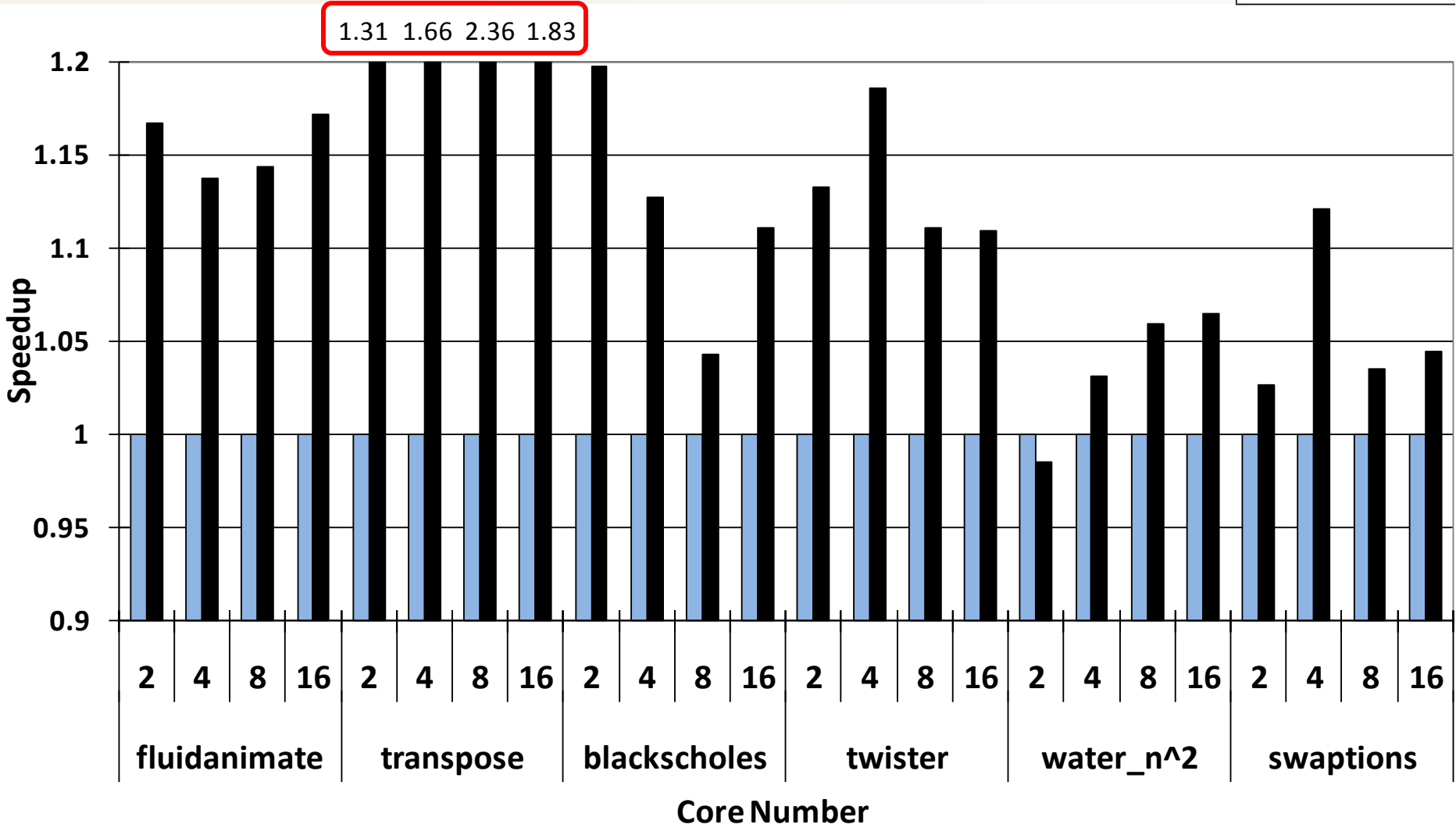
Experimental Setup

- 2 cores
 - Intel Core 2 Duo 6600 (2.4 Ghz)
- 4 cores
 - Intel Core 2 Quad Q6600 (2.4.Ghz)
- 8 cores
 - 2 Quad-core CPUs with SMT
 - Intel Xeon E5520 (2.26 Ghz)
- 16 cores (Logical)
 - 2 Quad-core CPUs with SMT and HyperThreading
 - Intel Xeon E5520 (2.26 Ghz)

Results



#thr=#core
Thread Tailor





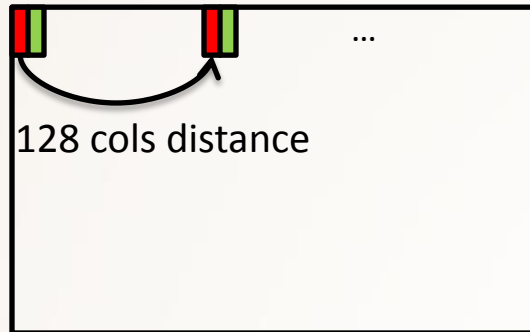
Result Analysis - Transpose

- Transpose $m * n$ matrix to $n * m$

$$\begin{pmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \end{pmatrix} \Rightarrow \begin{pmatrix} 1 & 4 \\ 2 & 5 \\ 3 & 6 \end{pmatrix}$$

- Parallel Transpose

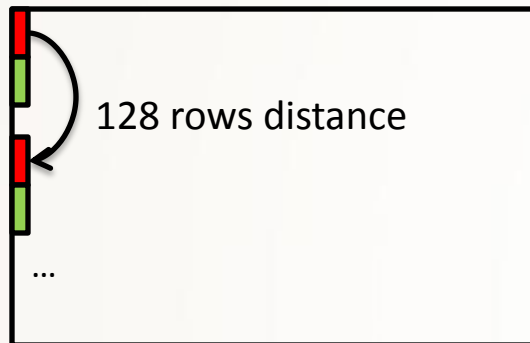
Input Matrix



Thread 1

Thread 2

Output Matrix





Result Analysis - Transpose

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Input Matrix
16K x 16K

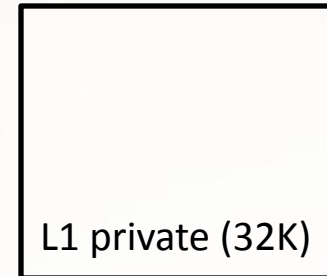


Output Matrix
16K x 16K



Intel Nehalem

Core 0

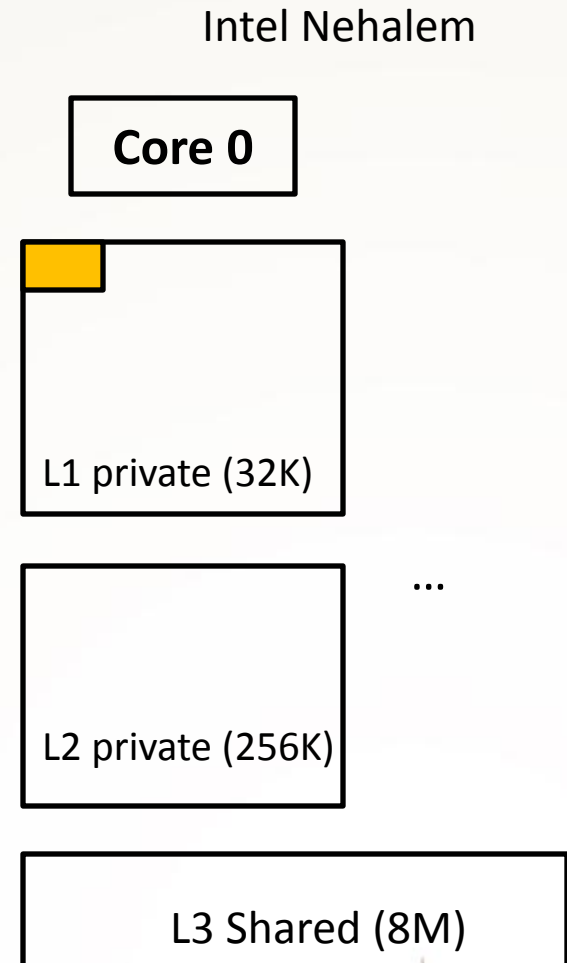
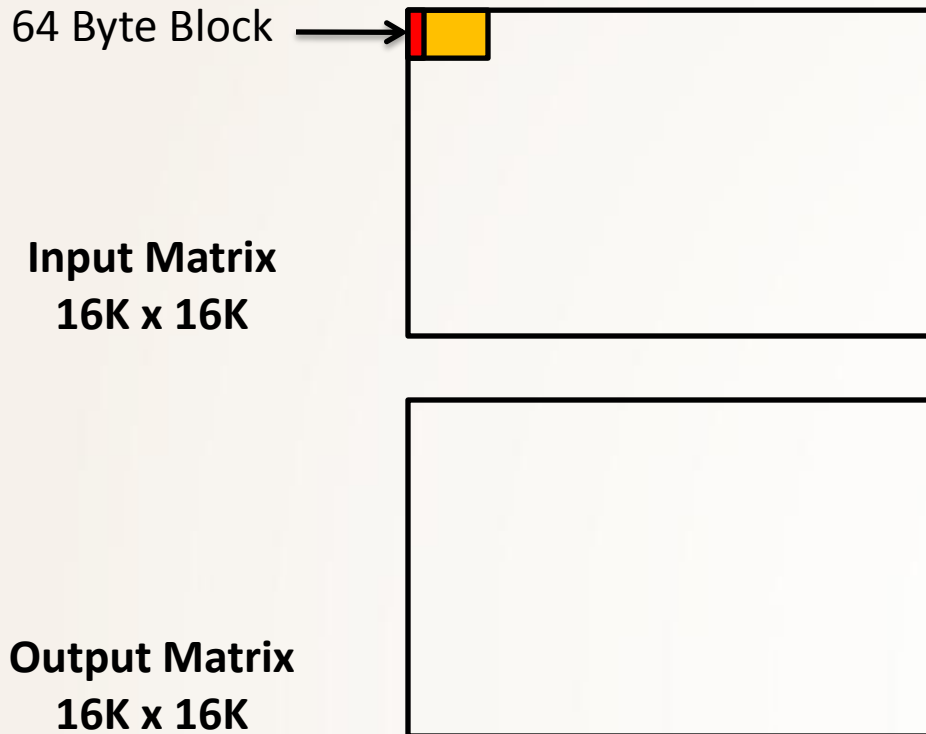




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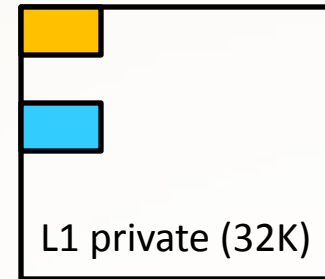


Output Matrix
16K x 16K



Intel Nehalem

Core 0



L1 private (32K)



L2 private (256K)



L3 Shared (8M)

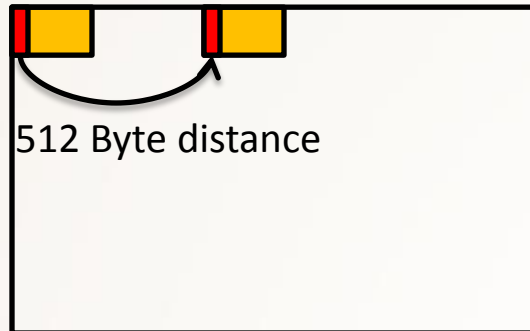


Result Analysis - Transpose

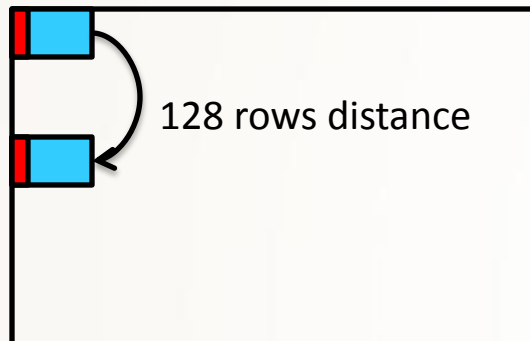
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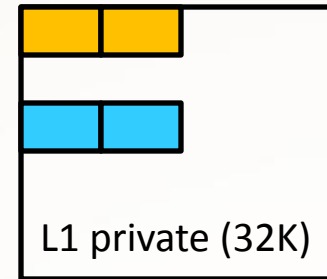


Output Matrix
16K x 16K



Intel Nehalem

Core 0



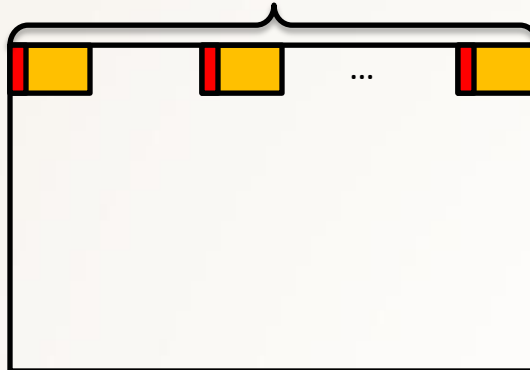


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iterates 128 times



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16K x 16K

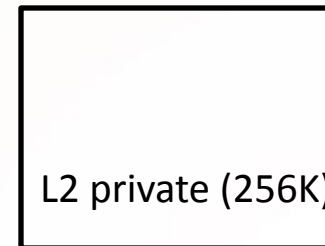
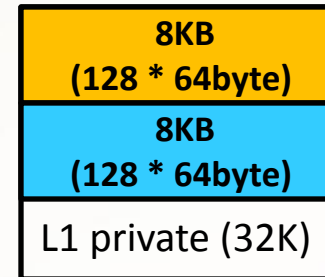
iterates 128 times



Output Matrix
16K x 16K

Intel Nehalem

Core 0



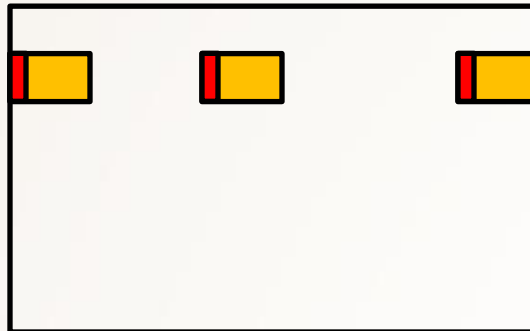


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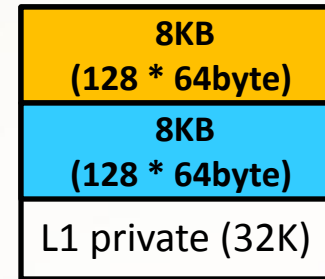


Output Matrix
16K x 16K



Intel Nehalem

Core 0



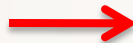
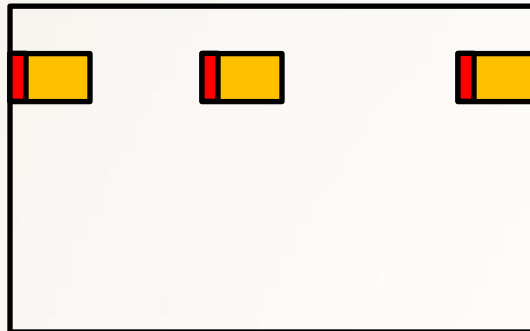


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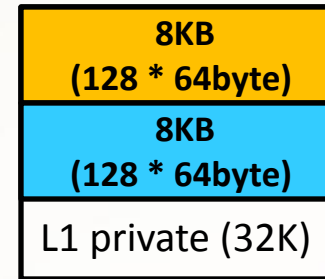


Output Matrix
16K x 16K



Intel Nehalem

Core 0



WRITE HIT!



...





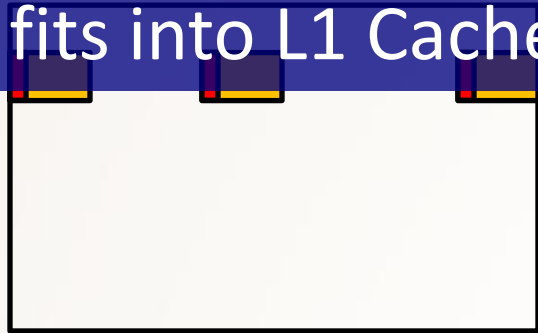
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Working-set fits into L1 Cache (No Capacity Miss!) WRITE HIT!

Input Matrix
16K x 16K

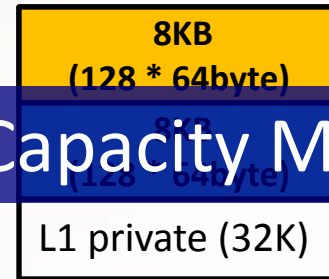


Output Matrix
16K x 16K



Intel Nehalem

Core 0





Summary

- Choosing Optimal Number of Threads is Hard
- Thread Tailor Ease the Pain
 - Graph Representation
 - Combine Threads at Runtime



Thank you